**Graphs 6**

**Weighted Edge List: Dijkstra's Algorithm**

The general problem is to find the shortest distances between two locations. There are several approaches:

* DFS (depth-first search): gives an answer, but not necessarily the shortest. O(V \* E)
* BFS (breadth-first search): give the shortest path in the sense of the smallest number of vertices or edges, but not necessarily the shortest weight (or distance) path. O(V \* E)
* brute force, try all paths: it works, but is slow. Floyd's is an example. O(V3)
* the *greedy* approach: always take the shortest outgoing path, without any long-term planning. Dijkstra's Algorithm is a classic example of a greedy algorithm. It is also an example of *dynamic programming*, meaning, to arrive at the answer by successive approximation. If we use a priority queue, the Big-O is O(V\*E\*log V)

Let’s develop Dijkstra’s algorithm. The internet has many discussions, videos, and demonstrations of Dijkstra's algorithm. Here some good demonstrations:

<http://optlab-server.sce.carleton.ca/POAnimations2007/DijkstrasAlgo.html>

<https://www.youtube.com/watch?v=UG7VmPWkJmA>

<https://www.youtube.com/watch?v=8Ls1RqHCOPw>

The general idea is to keep both edge-distances and a minimum distance from the source to every other vertex. As the vertices are processed, the distances back to the source are updated with shorter distances, if possible. The algorithm runs faster if the next vertex to be processed is always the smallest distance of the available vertexes (that's the priority queue part).

Look at this graph and visually calculate the shortest distance from A to B. \_\_\_\_

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
|  | A | B | C | D |
| start here | 0 | ∞ | ∞ | ∞ |
| after v=A | 0 | 9 | 3 | inf |
| after v=B |  |  |  |  |
| after v=C | 0 | 8 | 3 | 5 |
| after v=B |  |  |  |  |
| after v=D | 0 | 6 | 3 | 5 |
| after v=B | 0 | 6 | 3 | 5 |

9

1

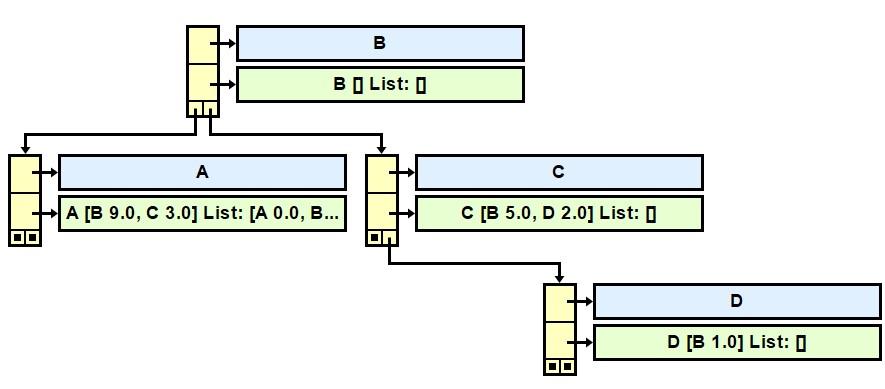
2

3

5

Filling in the table, calculate the shortest distance from A to every other vertex using Dijkstra's Algorithm. The source is A. The distances to the other vertices start at infinity. Take each neighbor of A and use their edge-distances to compute, if possible, a new shorter distance back to each already-processed vertex. Somehow, update the new distance for each vertex. Write the successive updates in this table. At the end, A's distance to itself is 0, A's distance to B is 6, A's distance to C is 3, and A's distance to D is 5. Genius.

What data structure can support all this? There are several. We will use this one:



First, the adjacency list is a Set of Neighbor,second, the List of PQelement or the Answers

I hope you see a

Map<String, wVertex> vertexMap = new TreeMap<String, wVertex>();

I hope you see, in A, the beginnings of this List: [A 0.0, B 6.0, C 3.0, D 5.0]. That's the answer. That's the minimum distances from A to every other vertex.

I hope you see that A (the blue box) maps to a weighted vertex wVertex (the green box). A wVertex is somewhat like a Vertex:

1. wVertex stores its own name.
2. wVertex has a Set that stores Neighbor objects.
3. wVertex has an arrayList of PQelement.
4. Here is the wVertexInterface:

public String getName();  
public ArrayList<PQelement> getAlDistanceToVertex();  
public PQelement getPQelement(wVertex v);  
public Double getDistanceToVertex(wVertex v);  
public void setDistanceToVertex(wVertex v, double m);  
public Set<Neighbor> getNeighbors();   
public void addAdjacent(wVertex v, double d);   
public String toString();

We need a new class, called the PQelement class.

1. PQelement stores a wVertex.
2. PQelement stores the distance to that wVertex.
3. PQelement implements a compareTo method, based on distance, to make the priority queue work.
4. PQelement has a toString method which you will need to implement.

We need a new class, called the Neighbor class. (Some programmers just use a Map<wVertex,Double>.)

1. Neighbor stores a wVertex object.
2. Neighbor stores the distance to that wVertex object.
3. Since we will put Neighbor objects into Sets, Neighbor needs to function in HashSets and TreeSets. Neighbors are "equal" if they have the same name. Make it so.

All these classes need constructors, accessors, modifiers, and toString.

Let's develop Dijkstra’s algorithm in more detail, this time including the priority queue:



|  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| A  0 |  |  |  |  |  |  |  |  |  |

9

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|  |  |
| --- | --- |
| A has a list of Neighbors (that is, neighboring wVertexes and edgeDistances).  In the wVertex A (the source) create an arrayList of PQelements, one PQelement for each wVertex, all set to infinity except the one for A which is set to 0  Add the PQelement corresponding to the source (A) to the pq | A🡪 [ A|0 B|∞ C|∞ D|∞]  pq🡪 [ A|0] |
| Loop while the pq still has elements |  |
| The nearest vertex, A, is removed from the pq. | A🡪 [ A|0 B|∞ C|∞ D|∞]  pq🡪 [] |
| A's Neighbor B is processed. Since 0+9 < ∞, B's distanceToVertex   is updated to 9. B is placed in the pq. | A🡪 [ A|0 B|9 C|∞ D|∞]  pq 🡪 [ B|9] |
| A's Neighbor C is processed. Since 0+3 < ∞, C's distanceToVertex   is updated to 3. C is placed in the pq. | A🡪 [ A|0 B|9 C|3 D|∞]  pq 🡪 [ B|9 C|3] |
| The nearest vertex, C, is automatically removed from the pq. | A🡪 [ A|0 B|9 C|3 D|∞]  pq 🡪 [ B|9] |
| C's Neighbor B is processed. Since 3+5 < 9, remove B and update its   distanceToVertex to 8. B is re-placed in the pq. | A🡪 [ A|0 B|8 C|3 D|∞]  pq 🡪 [ B|8] |
| C's Neighbor D is processed. Since 3+2 < ∞, remove D and updates its   distanceToVertex to 5. D is placed in the pq. | A🡪 [ A|0 B|8 C|3 D|5]  pq 🡪 [ B|8 D|5] |
| The nearest vertex, D, is automatically removed from the pq. | A🡪 [ A|0 B|8 C|3 D|5]  pq 🡪 [ B|8] |
| D's Neighbor B is processed. Since 5+1 < 8, remove B and update its   distanceToVertex to 6. B is re-placed in the pq. | A🡪 [ A|0 B|6 C|3 D|5]  pq 🡪 [ B|6] |
| The nearest vertex, B, is automatically removed from the pq. | A🡪 [ A|0 B|6 C|3 D|5]  pq 🡪 [] |
| B has no neighbors. The pq is now empty.  End Loop  The wVertex A is now storing in the array of PQelements [ A|0 B|6 C|3 D|5]  The minimum distance from A to A is 0, to B is 6, to C is 3, and to D is 5. |  |

Write the pseudocode for Dijkstra's Algorithm, aka minimumWeightPath.

private void minimumWeightPath(String vName) //Dijkstra's

{

}

The graph object, implementing an interface, is similar to AdjList, updated to use wVertex.

class AdjListWeighted implements AdjListWeightedInterface {  
 //we want our map to be ordered alphabetically by vertex name  
 private Map<String, wVertex> vertexMap = new TreeMap<String, wVertex>();

/\* constructor is not needed! \*/

public Set<wVertex> getVertices()

public Map<String, wVertex> getVertexMap()  
 public wVertex getVertex(String vName)   
 public void addVertex(String vName)

public void addEdge(String source, String target, double d)

public void minimumWeightPath(String vName); //Dijkstra's

public String toString();

**Assignment:** Finish Neighbor, PQelement, wVertex and AdjListWeighted. The driver is Dijkstra\_6\_Driver. Notice that we hard-code all inputs and test each class individually. You will turn in AdjListWeighted.

**Sample Run** (Dijkstra\_6\_Driver.java using AdjListWeighted)

Test the wVertex class  
get the names:  
 alpha  
 beta  
get the list of Neighbors:   
 [beta 5.0]  
 [alpha 3.0]  
toString() shows the name, the name and distance to the neighbor(s), and the list of PQelements:   
 alpha [beta 5.0] List: []  
 beta [alpha 3.0] List: []  
   
Testing Neighbor  
Neighbor's toString(): alpha 100.0  
Neighbor's toString(): alpha 101.0  
Neighbors are equal if the names are the same: 0  
   
Testing PQelement  
PQelement's toString(): alpha 100.0  
PQelement's toString(): beta 100.0  
Comparing two PQelements returns the difference in distance: 0  
   
Hard-coding vertices and neighbors with weights.  
Get the vertex by name "A": A [B 9.0, C 3.0] List: []  
Get the vertices: [A [B 9.0, C 3.0] List: [], B [] List: [], C [B 5.0, D 2.0] List: [], D [B 1.0] List: []]  
Get the map: {A=A [B 9.0, C 3.0] List: [], B=B [] List: [], C=C [B 5.0, D 2.0] List: [], D=D [B 1.0] List: []}  
The whole graph:  
A [B 9.0, C 3.0] List: []  
B [] List: []  
C [B 5.0, D 2.0] List: []  
D [B 1.0] List: []  
   
Dijkstra's Algorithm!  
Enter source: A  
   
After processing, the entire graph is:  
A [B 9.0, C 3.0] List: [A 0.0, B 6.0, C 3.0, D 5.0]  
B [] List: []  
C [B 5.0, D 2.0] List: []  
D [B 1.0] List: []  
   
State of the source vertex: A  
A [B 9.0, C 3.0] List: [A 0.0, B 6.0, C 3.0, D 5.0]  
   
The source A knows the distance to each target:  
Distance to A: 0.0  
Distance to B: 6.0  
Distance to C: 3.0  
Distance to D: 5.0